

## Shoe for Sale

### Nude High Cut Shoe

- 10cm Natural Leather Stacked Heel.
- Multiple Strap Detail With Gold Western Buckles & Functional Zip.
- Leather Lining & Sock
- Unique

The original price of a pair was **\$500**.  
Now the last piece for **\$5** only.



# Combinatorial auctions

- A combinatorial auction sells multiple objects simultaneously.
- Bidders can place bids on combinations of items in **bundles**.
- Bidder's valuation on a bundle may be different from the sum of the valuations of all the items in the bundle.
  - **Complementary**: the value of a combination of items is worth more than the sum of the values of the separate items.
  - **Substitutable**: the value of a combination of items is less than the sum of the values of the separate items.

# The model of combinatorial auctions

$E = (N \cup \{0\}, X, \{v_i\}_{i \in N})$  is a combinatorial auction if

- $N = \{1, 2, \dots, n\}$  is the set of buyers
- 0 represents the seller
- $X$  is the set of items
- $v_i : 2^X \rightarrow \mathbb{Z}^+$  the buyer  $i$ 's value function

## Example:

$N = \{1, 2\}, X = \{a, b\}$ .

$v_1(\emptyset) = 0, v_1(\{a\}) = v_1(\{b\}) = v_1(\{a, b\}) = 1.$

$v_2(\emptyset) = 0, v_2(\{a\}) = v_2(\{b\}) = 1, v_2(\{a, b\}) = 3.$

**Question:** How to allocate the items to the buyers so that each item goes to the buyer who gives it the highest value?

# Efficient allocations

- **Allocation:**  $\pi : N \cup \{0\} \rightarrow 2^X$  such that
  - $\pi(i) \cap \pi(j) = \emptyset$  for any  $i \neq j$ .
  - $\bigcup_{i \in N \cup \{0\}} \pi(i) = X$ .

which allocate all the items to the buyers, each buyer can have a bundle but one item can only be allocated to at most one buyer.

- **Efficient allocation**  $\pi^*$ :  $\pi^*(0) = \emptyset$  and for every allocation  $\pi$  of  $X$ ,

$$\sum_{i \in N} v_i(\pi^*(i)) \geq \sum_{i \in N} v_i(\pi(i))$$

**Question:** How to find an efficient allocation?

# Walrasian equilibria

- **Price vector  $\mathbf{p}$** : assign a non-negative real number to each item in  $X$ .
- **Demand correspondence**:

$$D_i(\mathbf{p}) = \arg \max_{A \subseteq X} (V_i(A) - \sum_{a \in A} p_a)$$

representing all the bundles that give  $i$  the highest utility based on the current market price. For instance, if  $\mathbf{p} = (0.5, 0.5)$ ,

$$D_1(\mathbf{p}) = \{\{a\}, \{b\}\}. \quad D_2(\mathbf{p}) = \{\{a, b\}\}$$

- **Walrasian equilibrium  $(\mathbf{p}, \pi)$** :  $\mathbf{p}$  is a price vector and  $\pi$  is an allocation of  $X$  such that  $\pi(0) = \emptyset$  and  $\pi(i) \in D_i(\mathbf{p})$  for all  $i \in N$ .
- Any Walrasian equilibrium determines an efficient allocation.

**Question:** How to find a Walrasian equilibrium?